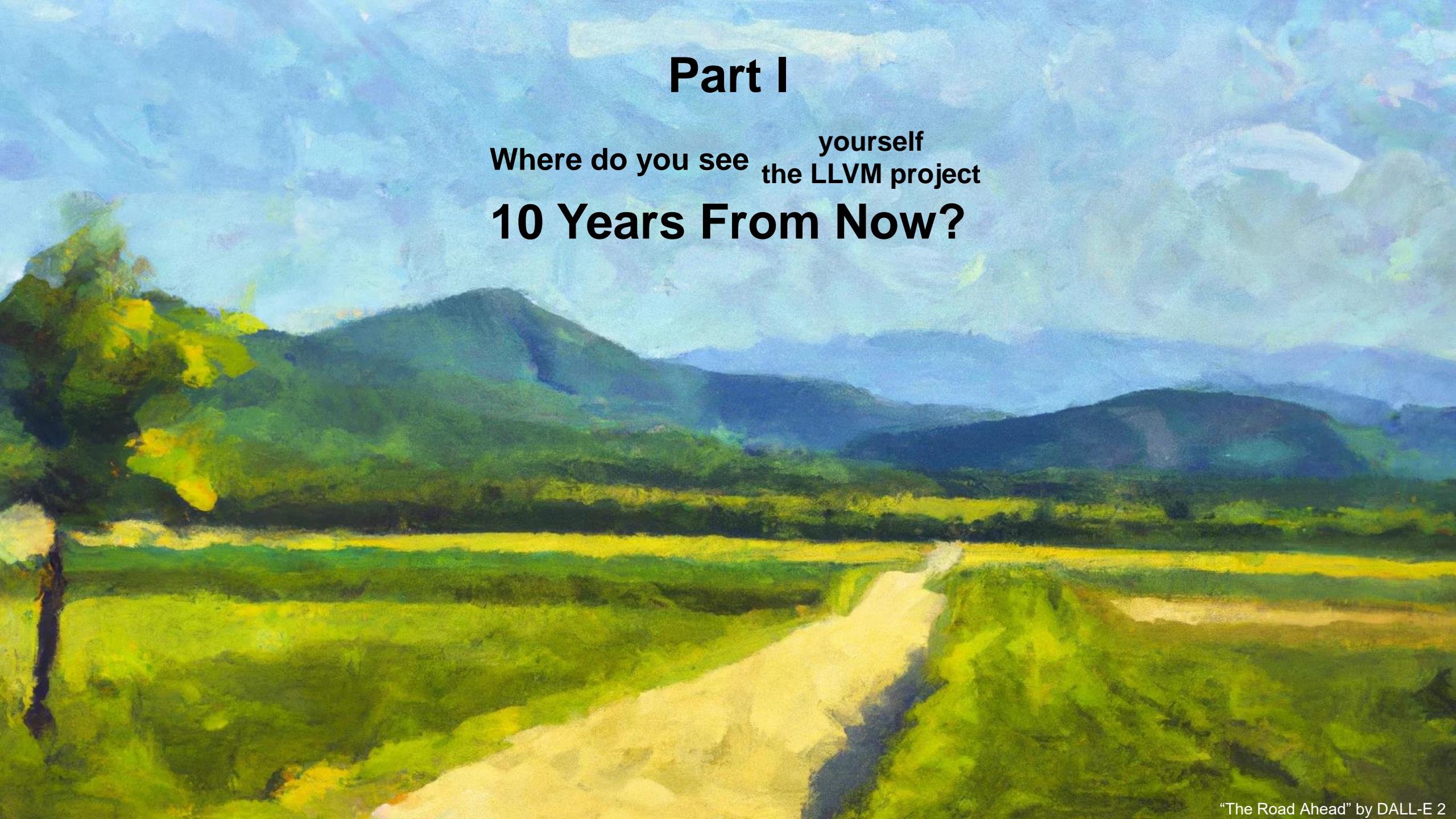




Paths towards unifying LLVM and MLIR

Nicolai Hähnle



Part I

Where do you see ^{yourself}
the LLVM project

10 Years From Now?

“The LLVM project”

The screenshot shows the LLVM GitHub repository page. The top navigation bar includes the repository name "Ivm / Ivm-project" (Public), "Edit Pins", "Unwatch 565", "Fork 6.3k", and "Star 16.6k". Below the navigation is a tab bar with "Code" (selected) and "Issues". The main content area features a large, bold title "Communication" followed by a horizontal line (~). Below this is another large, bold title "Intermediate Representation". To the left, there is a sidebar with a tree view of repository branches: bolt, clang-tools-extra, clang, cmake, compiler-rt, and cross-project-tests. The main content area displays a list of recent commits:

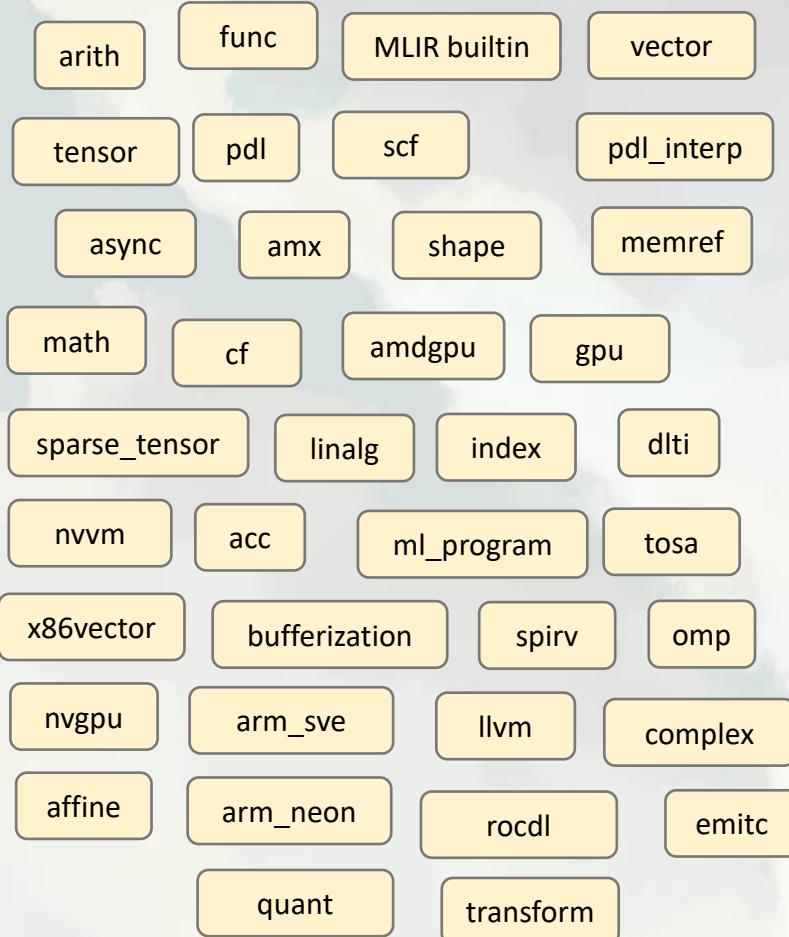
Commit	Message	Time Ago
bolt	[clangd] Add scopea enum constants to all se...	2 days ago
clang-tools-extra	[clang] Remove an incorrect assert	1 hour ago
clang	Harmonize cmake_policy() across standalone bu...	7 days ago
cmake	[sanitizer] Fix vfork interception on loongarch64	2 days ago
compiler-rt	[dexter-tests] Add attribute optnone to main fu...	9 days ago
cross-project-tests		

To the right of the commit list, there is a sidebar with links: "Submit your patches at <http://reviews.llvm.org>", "Ivm.org", "Readme", "View license", and "Security policy".

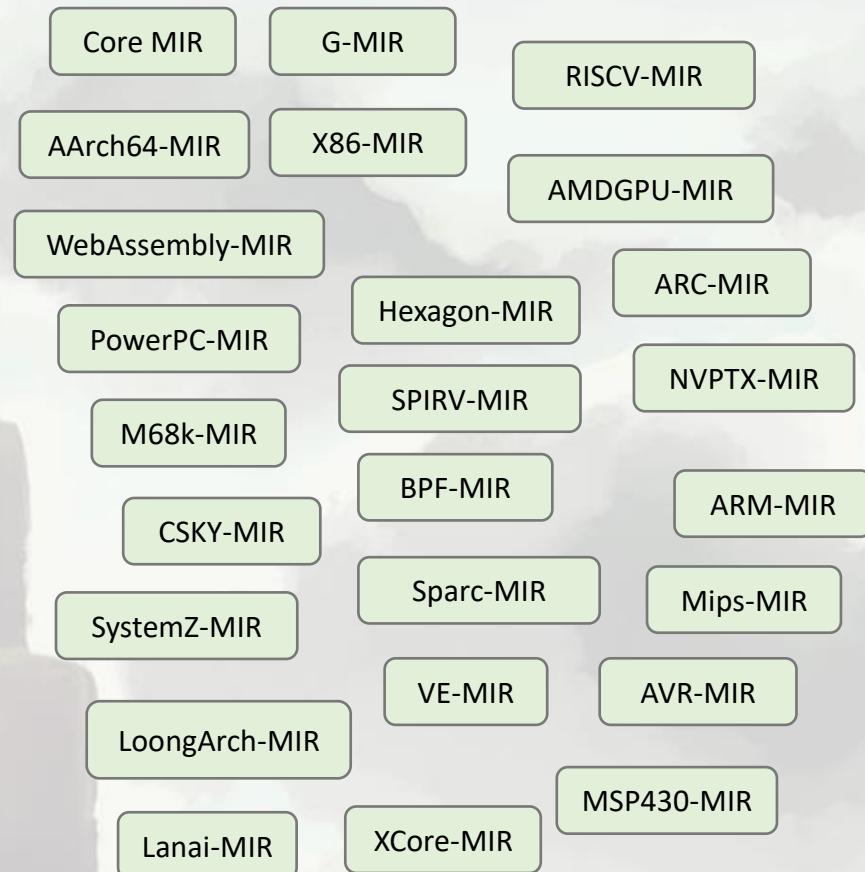


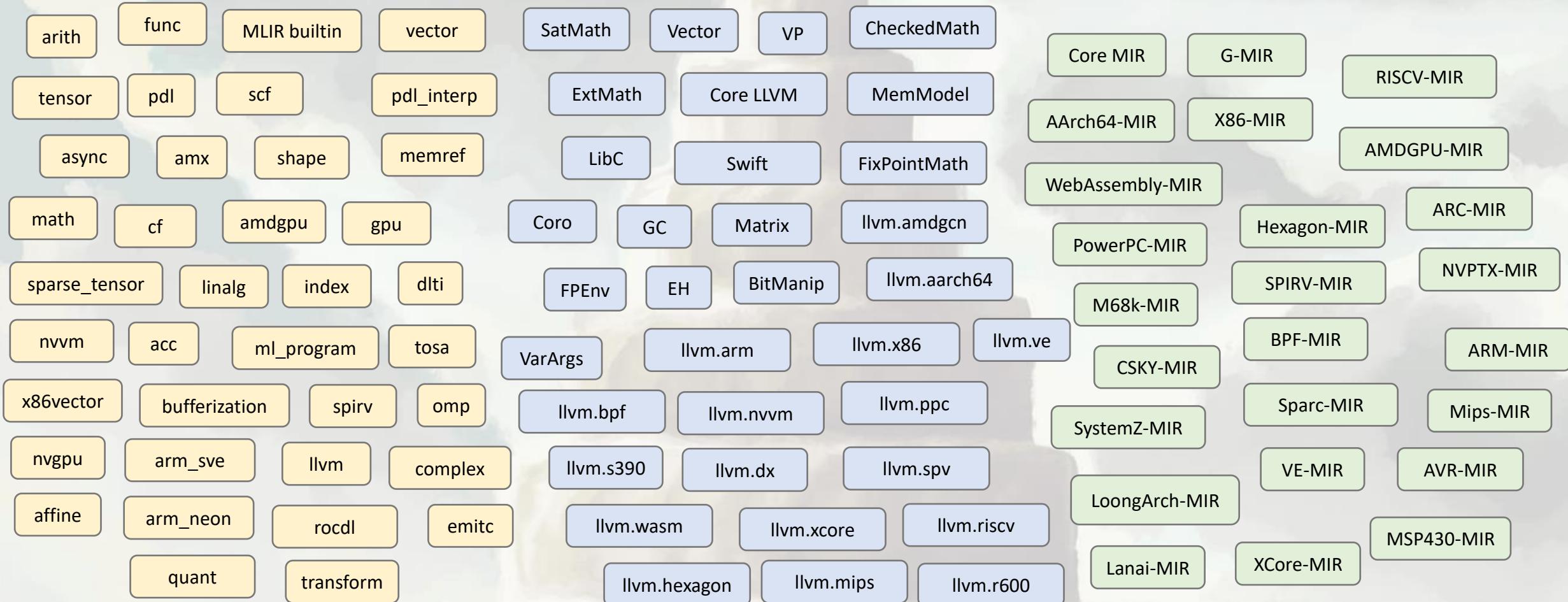
LLVM IR





LLVM IR





What is “LLVM IR”?

```
%34 = and i32 %23, %24
%.not = icmp ne i32 %34, -1
%35 = zext i1 %.not to i32
%.not19 = icmp eq i32 %28, 0
br i1 %.not19, label %53, label %36

36:
br i1 %.not, label %37, label %44

37:
%38 = lshr i32 %33, 2
%39 = zext i32 %38 to i64
```

```
000001a0: 001f 6c40 9aff ffff ff1f 0009 e803 800e ..l@.....
000001b0: 0000 0000 4918 0000 0600 0000 1382 6082 ....I.....
000001c0: 200c 1302 62c2 5018 c784 0249 9485 9920 ...b.P....I...
000001d0: 3407 0000 1a21 4c0e f3e0 857e b7cb 2f37 4....!L....~/7
000001e0: 3a5c 46bb d9e5 57b8 4d7e 89eb e7f9 6b3e :\F...W.M~....k>
000001f0: 8eb5 6638 9bad 7536 e363 b156 190f 7fcf ..f8..u6.c.V....
00000200: e9f3 97bc 4cae 8fcf 6417 9b3d a402 9e00 ....L...d..=....
00000210: 1700 2000 0000 0000 0000 0000 7880 21d5 .. .....x!.!
00000220: f315 4000 0800 0000 0000 0000 0f30 ..@.....0
00000230: a472 8306 0102 6000 0000 0000 0000 0000 ..r.....
00000240: 7880 21d5 1e30 0a10 0003 0000 0000 0000 x.!..0.....|
```

LLVM Language Reference Manual ↗

- Abstract
- Intro

Dialect

- Abstract
- Calling Conventions
- Visibility Styles
- DLL Storage Classes
- Thread Local Storage Models
- Runtime Preemption Specifiers

Substrate

```
Value *isEvenI = m_builder->CreateICmpEQ(
    m_builder->CreateSMod(m_builder->CreateFPToSI(xyChromaInfo.coordI, m_
    m_builder->CreateFPExt(0.0f, 0.0f), m_builder->CreateFPExt(1.0f, 1.0f)),
    m_builder->CreateFPMul(m_builder->CreateFPExt(0.0f, 0.0f), m_builder->CreateFPExt(1.0f, 1.0f)));
Value *isEvenJ = m_builder->CreateICmpEQ(
    m_builder->CreateSMod(m_builder->CreateFPToSI(xyChromaInfo.coordJ, m_
    m_builder->CreateFPExt(0.0f, 0.0f), m_builder->CreateFPExt(1.0f, 1.0f)),
    m_builder->CreateFPMul(m_builder->CreateFPExt(0.0f, 0.0f), m_builder->CreateFPExt(1.0f, 1.0f)));
Value *subCoordI = m_builder->CreateFPMul(
    Intrinsic::floor, m_builder->CreateFPExt(xyChromaInfo.coordI, ConstantFP(0.0f)),
    Value *subCoordJ = m_builder->CreateUnaryIntrinsic(
        Intrinsic::floor, m_builder->CreateFDiv(xyChromaInfo.coordJ, ConstantFP(0.0f))));
```

Dialect

Instruction set
Types
Semantics

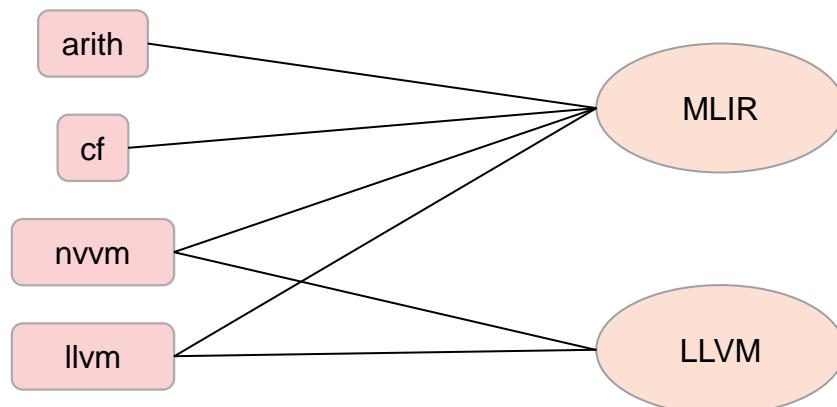
E.g.: LangRef.rst, mlir/Dialect.td

Substrate

Set of C++ classes to represent and manipulate code in one or more dialects

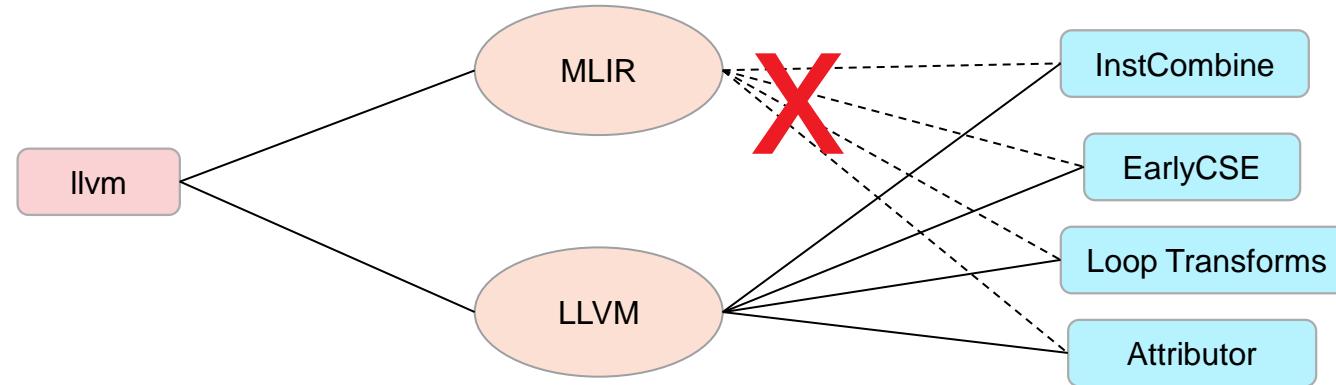
E.g.: llvm::Instruction, llvm::Value, llvm::BasicBlock, ...

N:M relationship



Substrates Matter

- Code artefacts (transforms, analyses, helpers, ...) are written against a specific substrate
- Can represent llvm in MLIR, but can't run InstCombine or loop transforms on it

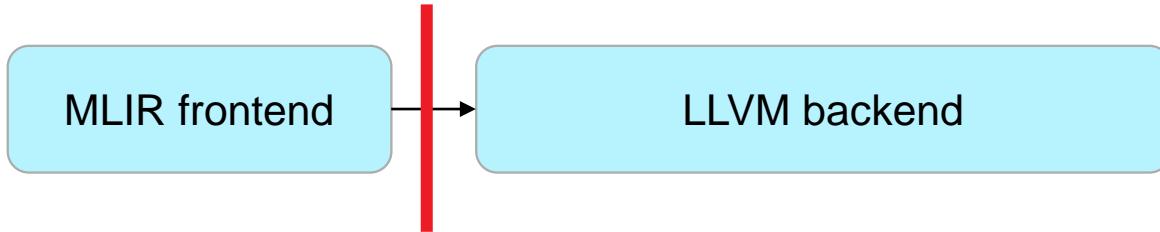


Substrates Matter

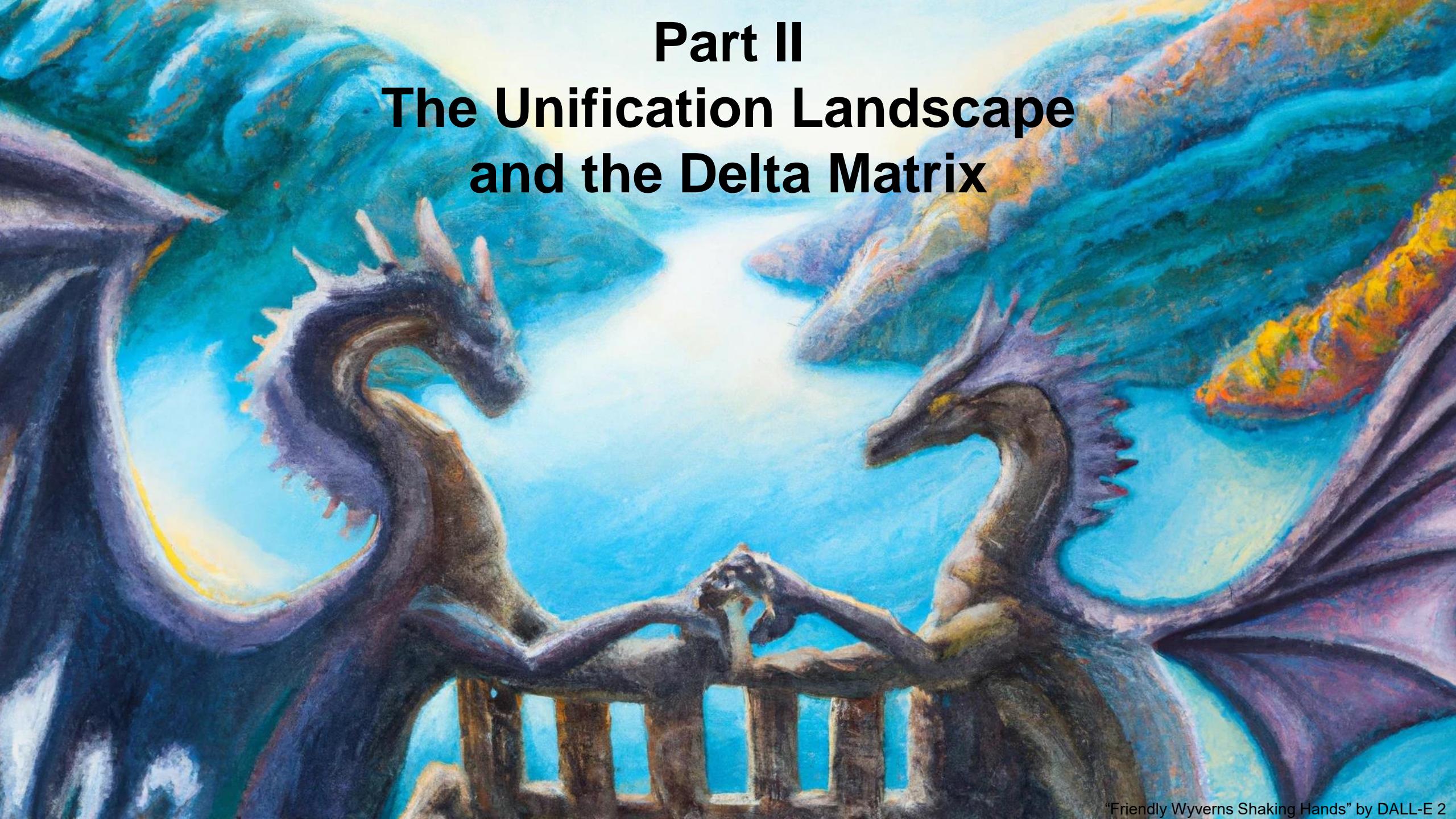
```
%2535 = call i32 @llvm.amdgcn.uffe.i32(i32 %.4.vec.extract1487, i32 16, i32 16)
%2536 = select i1 false, i32 %.4.vec.extract1487, i32 %2535
%2537 = select i1 false, i32 0, i32 %2536
%2538 = uitofp i32 %2537 to float
%2539 = call i32 @llvm.amdgcn.uffe.i32(i32 %.4.vec.extract1487, i32 0, i32 16)
%2540 = select i1 false, i32 %.4.vec.extract1487, i32 %2539
%2541 = select i1 false, i32 0, i32 %2540
%2542 = uitofp i32 %2541 to float
%2543 = insertelement <4 x float> undef, float %2530, i64 0
%2544 = insertelement <4 x float> %2543, float %2534, i64 1
%2545 = insertelement <4 x float> %2544, float %2538, i64 2
%2546 = insertelement <4 x float> %2545, float %2542, i64 3
%scale.i152 = fmul reassoc nnan nsz arcp contract afn <4 x float> %2546, <float 0x3F00002000000000, float 0x3F00002000000000, float 0x3F00002000000000>
%2547 = call i8 addrspace(4)* @lgc.descriptor.table.addr(i32 1, i32 1, i32 0, i32 23, i32 -1) #4
%2548 = getelementptr i8, i8 addrspace(4)* %2547, i32 512
%2549 = bitcast i8 addrspace(4)* %2548 to <4 x i32> addrspace(4)*
%2550 = call i8 addrspace(4)* @lgc.descriptor.table.addr(i32 2, i32 2, i32 0, i32 23, i32 -1) #4
%2551 = getelementptr i8, i8 addrspace(4)* %2550, i32 512
%2552 = bitcast i8 addrspace(4)* %2551 to <4 x i32> addrspace(4)*
%2553 = load <4 x i32>, <4 x i32> addrspace(4)* %2552, align 16
%2554 = load <8 x i32>, <8 x i32> addrspace(4)* %2549, align 32
%2555 = shufflevector <4 x float> %scale.i152, <4 x float> poison, <2 x i32> <i32 0, i32 1>
%2556 = fmul reassoc nnan nsz arcp contract afn <2 x float> %2492, %2555
%2557 = shufflevector <4 x float> %scale.i152, <4 x float> poison, <2 x i32> <i32 2, i32 3>
%2558 = fadd reassoc nnan nsz arcp contract afn <2 x float> %2556, %2557
%2559 = extractelement <2 x float> %2558, i64 0
%2560 = extractelement <2 x float> %2558, i64 1
%2561 = call reassoc nnan nsz arcp contract afn <4 x float> @llvm.amdgcn.image.sample.l.2d.v4f32.f32(i32 15, float %2559, float %2560, float 0.000000)
%2562 = extractelement <4 x float> %2561, i64 0
%2563 = fmul reassoc nnan nsz arcp contract afn float %2562, %2562
%2564 = icmp sg i32 %2571, 1
```

Benefits of a Unified Substrate

- Interoperability of code
- Eliminate a hard pass-ordering boundary



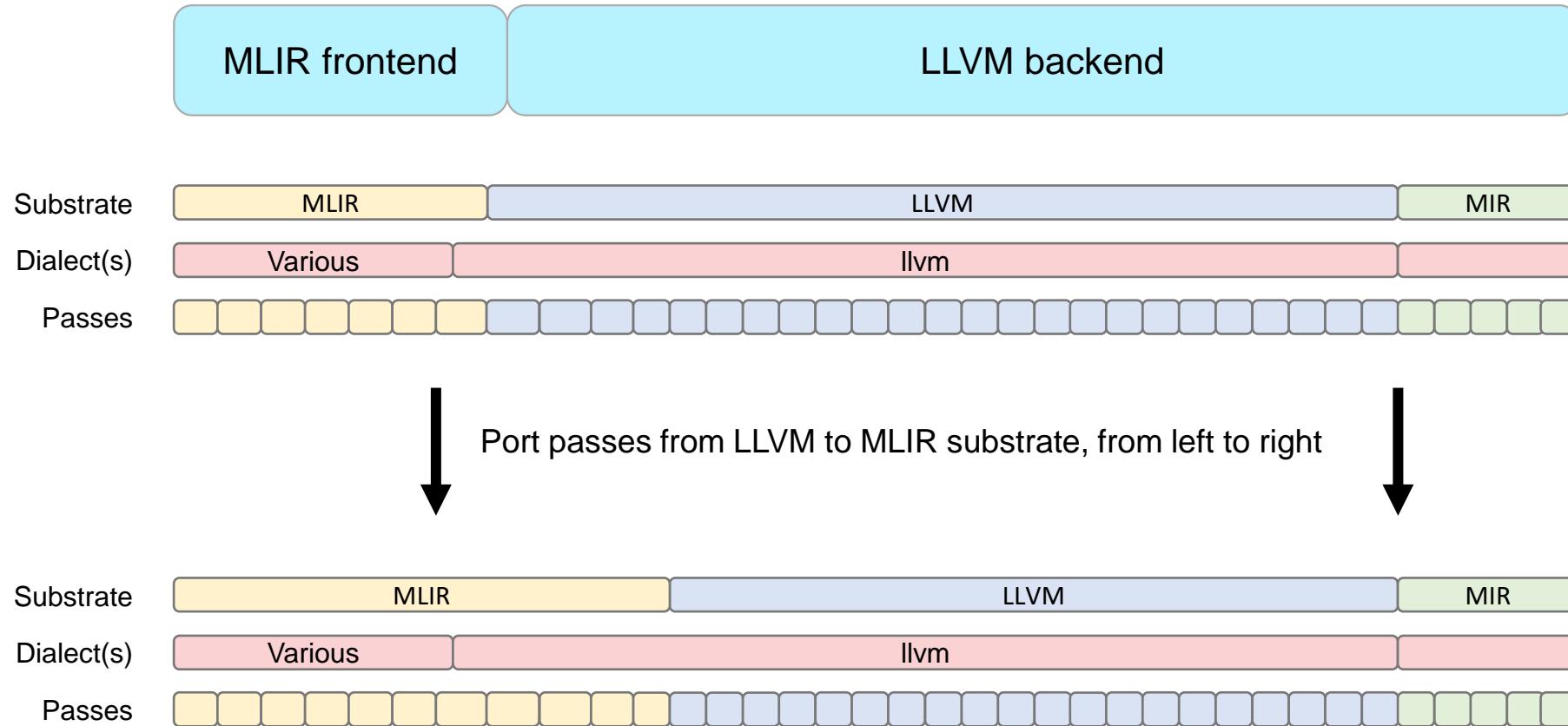
- Accessibility for developers
 - Shared knowledge and shared “language” within the LLVM project
 - Reduced friction for “full (compiler) stack developers”



Part II

The Unification Landscape and the Delta Matrix

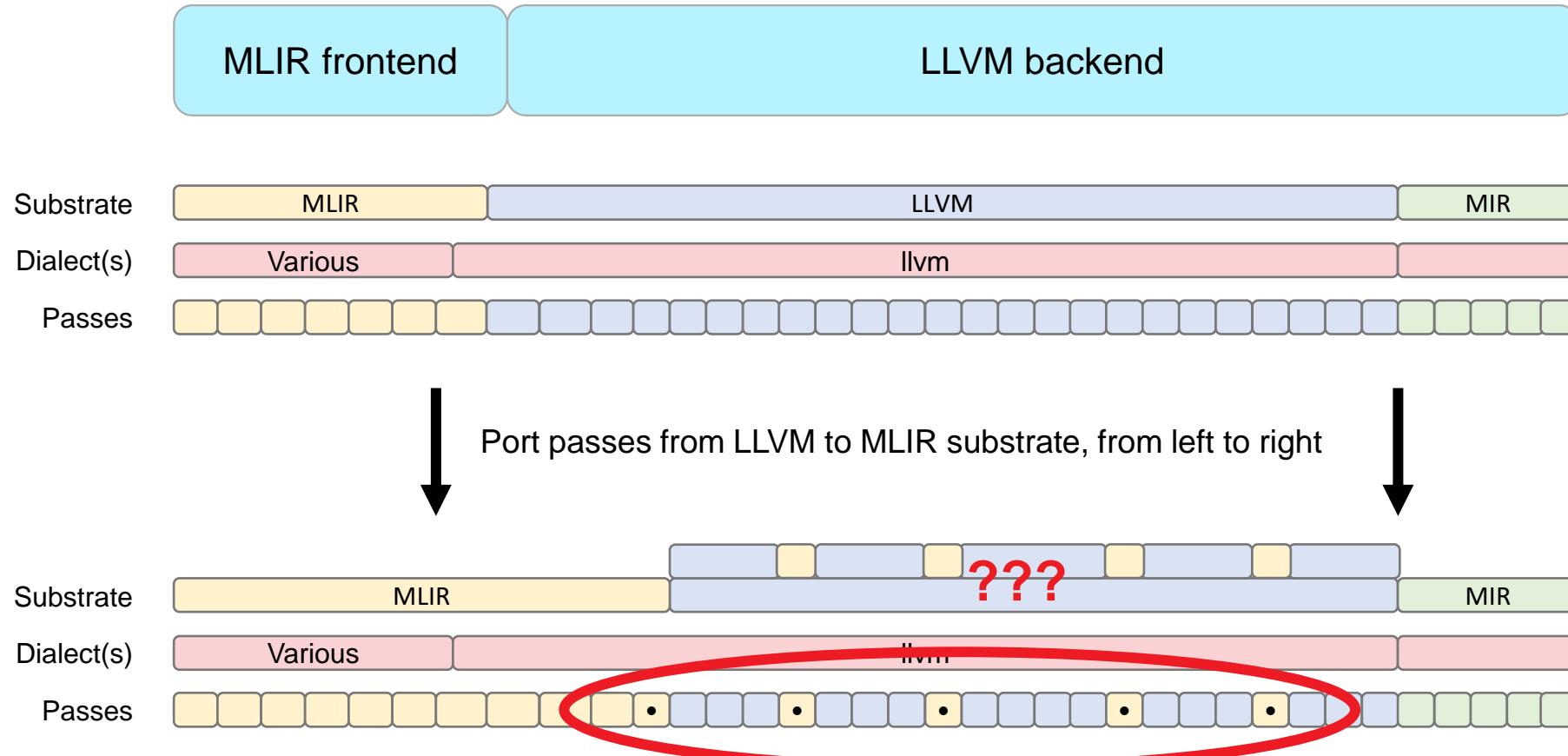
Thought Experiment: Port LLVM Code to the MLIR Substrate



Reality Check: The Optimization Pass Pipeline

```
$ opt -O3 -print-pipeline-passes -S /dev/null
verify,annotation2metadata,forceattrs,inferattrs,coro-early,function<eager-inv>(lower-expect simplifycfg bonus-inst-threshold=1;
no-forward-switch-cond;no-switch-range-to-icmp;no-switch-to-lookup;keep-loops;no-hoist-common-insts;no-sink-common-insts>,sroa,
early-cse<>,callsite-splitting),openmp-opt,ipscce,called-value-propagation,globalopt,function(mem2reg),function<eager-inv>
(instcombine simplifycfg bonus-inst-threshold=1;no-forward-switch-cond;switch-range-to-icmp;no-switch-to-lookup;keep-loops;
no-hoist-common-insts;no-sink-common-insts>),require<globals-aa>,function(invalidate<aa>),require<profile-summary>,cgscce(devirt<4>
(inline<only-mandatory>,inline,function-attrs,argpromotion,openmp-opt-cgscce, function<eager-inv>(sroa,early-cse<memssa>,
speculative-execution,jump-threading,correlated-propagation simplifycfg bonus-inst-threshold=1;no-forward-switch-cond;
switch-range-to-icmp;no-switch-to-lookup;keep-loops;no-hoist-common-insts;no-sink-common-insts>,instcombine,aggressive-instcombine,
libcalls-shrinkwrap,tailcallelim simplifycfg bonus-inst-threshold=1;no-forward-switch-cond;switch-range-to-icmp;no-switch-to-lookup;
keep-loops;no-hoist-common-insts;no-sink-common-insts>,reassociate,require<opt-remark-emit>,loop-msa(loop-instsimplify,
loop-simplify,licm<no-allowspeculation>,loop-rotate,licm<allowspeculation>,simple-loop-unswitch<nontrivial;trivial>),
simplifycfg bonus-inst-threshold=1;no-forward-switch-cond;switch-range-to-icmp;no-switch-to-lookup;keep-loops;no-hoist-common-insts;
no-sink-common-insts>,instcombine,loop(loop-idiom,indvars,loop-deletion,loop-unroll-full),sroa,mldst-motion<no-split-footer-bb>,
avn<>.scce,instcombine,jump-threading,correlated-propagation,adce,memcpyopt,dse,loop-msa(licm<allowspeculation>),coro-elide,
simplifycfg bonus-inst-threshold=1;no-forward-switch-cond;switch-range-to-icmp;no-switch-to-lookup;keep-loops;hoist-common-insts;
sink-common-insts>,instcombine),coro-split)),deadargelim,coro-cleanups,globalopt,globaldce,elim-avail-extern,rpo-function-attrs,
recompute-globalssaa,function<eager-inv>(float2int,lower-constant-intrinsics,loop(loop-rotate,loop-deletion),loop-distribute,
inject-tli-mappings,loop-vectorize<no-interleave-forced-only;no-vectorize-forced-only;>,loop-load-elim,instcombine,
simplifycfg bonus-inst-threshold=1;forward-switch-cond;switch-range-to-icmp;switch-to-lookup;no-keep-loops;hoist-common-insts;
sink-common-insts>,slp-vectorizer,vector-combine,instcombine,loop-unroll<O3>,transform-warning,instcombine,require<opt-remark-emit>,
loop-msa(licm<allowspeculation>),alignment-from-assumptions,loop-sink,instsimplify,div-rem-pairs,tailcallelim,
simplifycfg bonus-inst-threshold=1;no-forward-switch-cond;switch-range-to-icmp;no-switch-to-lookup;keep-loops;no-hoist-common-insts;
no-sink-common-insts>),globaldce,constmerge,cg-profile,rel-lookup-table-converter,function(annotation-remarks),verify,print
```

Thought Experiment: Port LLVM Code to the MLIR Substrate



- Many substrate round-trips (compile-time cost!) or code duplication (maintenance cost!)

Thought Experiment: Make Code Generic over the Substrate

- Generic code already exists. For example:

```
template <typename NodeT, bool IsPostDom>
class DominatorTreeBase {
```

Expand use of this technique. For example:

```
template <typename BlockT, typename InstructionT>
void hoistAllInstructionsInto(BlockT *DomBlock, InstructionT *InsertPt,
                             BlockT *BB);
```

Eventually:

```
namespace llvm {
using BasicBlock = mlir::Block;
using Instruction = mlir::Operation;
// ...
}
```

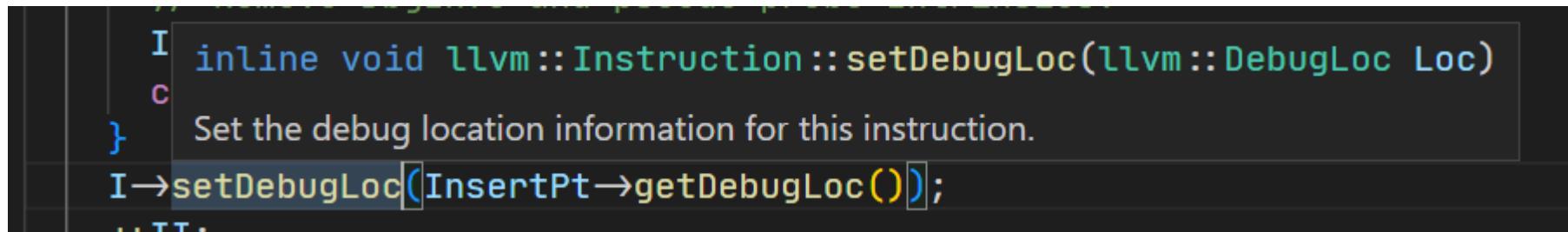


```
template <typename BlockT, typename InstructionT>
void llvm::hoistAllInstructionsInto(BlockT *DomBlock, InstructionT *InsertPt,
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
// (...)

for (typename BlockT::iterator II = BB->begin(), IE = BB->end(); II != IE;) {
    InstructionT *I = &*II;
    I->drop.UndefImplviaAttrsAndUnknownMetadata();
    if (I->isUsedByMetadata())
        dropDebugUsers(*I);
    if (I->isDebugOrPseudoInst()) {
        // Remove DbgInfo and pseudo probe Intrinsics.
        II = I->eraseFromParent();
        continue;
    }
    I->setDebugLoc(InsertPt->getDebugLoc());
    ++II;
}
DomBlock->getInstList().splice(InsertPt->getIterator(), BB->getInstList().
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
BB->begin(),
BB->getTerminator()>getIterator());
}
```

Tooling for Templates is Weak

- Error messages
- Compile times
- Language servers



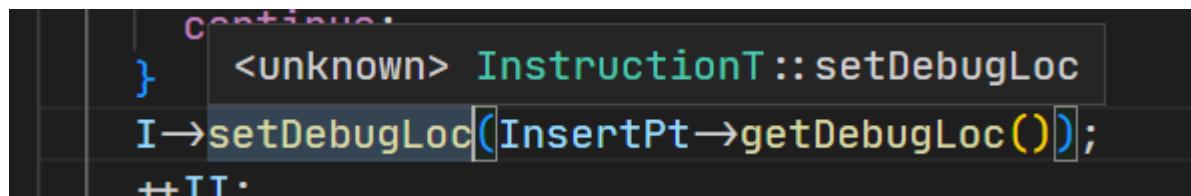
A screenshot of an IDE showing a tooltip for a template function. The code snippet is:

```
I inline void llvm::Instruction::setDebugLoc(llvm::DebugLoc Loc)
}
Set the debug location information for this instruction.

I→setDebugLoc(InsertPt→getDebugLoc());
```

The tooltip contains the documentation: "Set the debug location information for this instruction."

VS.



A screenshot of an IDE showing a tooltip for a template function. The code snippet is identical to the one above:

```
continuous
} <unknown> InstructionT::setDebugLoc
I→setDebugLoc(InsertPt→getDebugLoc());
#TT·
```

The tooltip is empty, showing only the placeholder text "<unknown>".

The Delta Matrix

Clang										
Flang										
SimplifyCFG										
InstCombine										
Attributor										
Coroutines										
AddressSanitizer										
Alias analysis										
Target codegen										
Assembly printer & parser										
Bitcode writer & reader										
IR linker										
lldb										
JIT infrastructure										
llvm-reduce										
...										

The Delta Matrix

	Constants	Multiple defined values	Phi nodes vs. BB args	Builder interface	Regions	Attributes	Debug info	Metadata	...
Clang									
Flang									
SimplifyCFG									
InstCombine									
Attributor									
Coroutines									
AddressSanitizer									
Alias analysis									
Target codegen									
Assembly printer & parser									
Bitcode writer & reader									
IR linker									
lldb									
llvm-reduce									
lli									
...									

llvm::Value



Part III

Column Refactorings

Example 1: Dialects in LLVM

1. First-class support for extended instructions and types, with near-MLIR-compatible definitions
2. Cleanup of “the LLVM dialect” into constituent dialects
3. Extensibility for frontends, including external frontends

First-class support for extended instructions and types

- Intrinsics are second-class citizens in textual IR

```
%x = add i32 %y, %z
```

vs.

```
%x = call i32 @llvm.smin.i32(i32 %y, i32 %z)
```

- Intrinsics are second-class citizens in C++

```
if (auto *Store = dyn_cast<StoreInst>(I)) {  
    use(Store->getPointerOperand());  
}
```

vs.

```
if (auto *Intr = dyn_cast<IntrinsicInst>(I)) {  
    if (Intr->getIntrinsicID() == Intrinsic::prefetch) {  
        use(Intr->getArgOperand(0));  
    }  
}
```

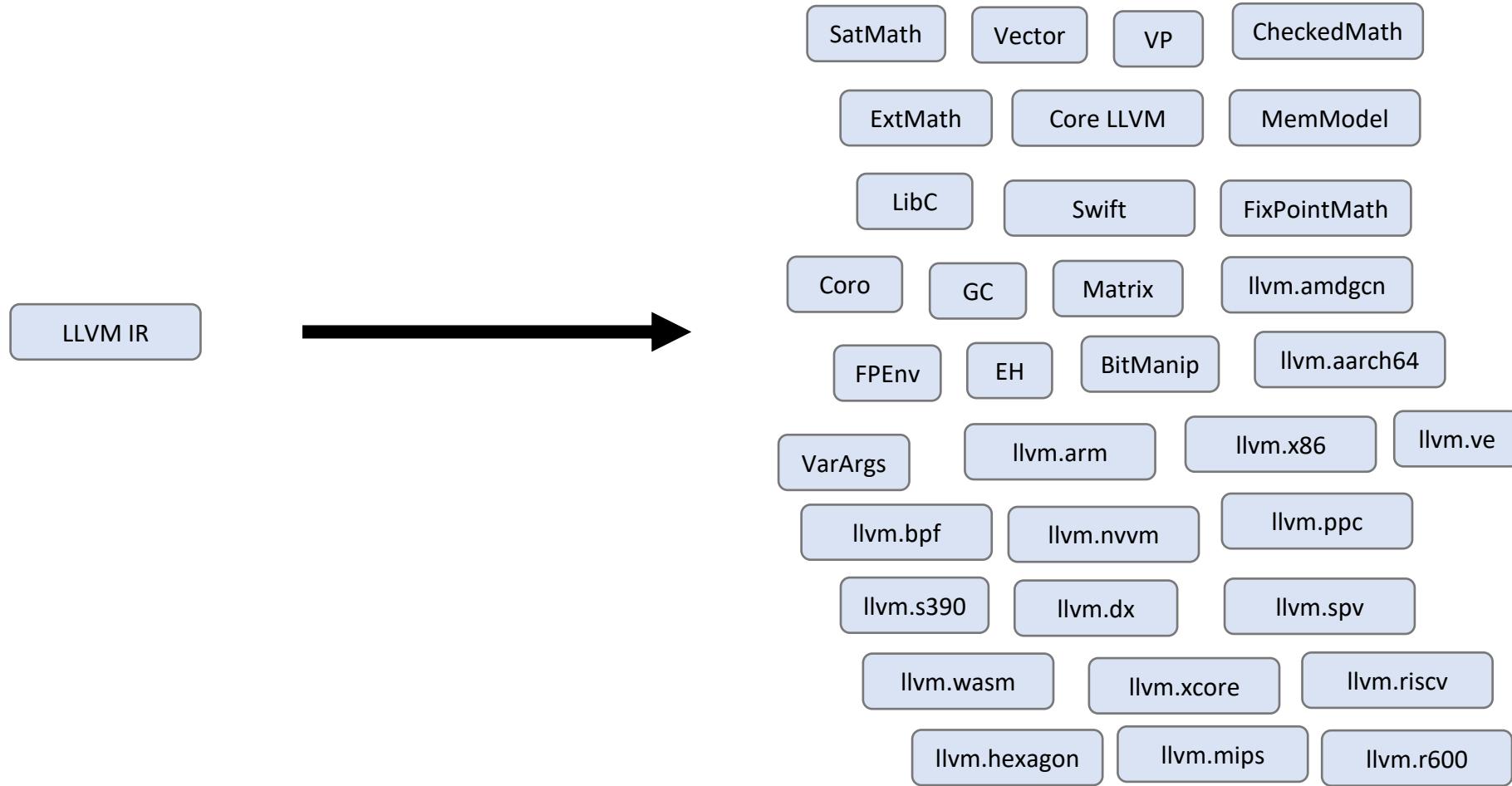
- Intrinsics are compile-time and -space inefficient

```
store volatile i32 %x, ptr %p
```

vs.

```
call void @llvm.memset.p0.i64(ptr %p, i8 %y, i64 4, i1 true)
```

Cleanup of “the LLVM dialect”



Extensibility

Today:

```
call void @llvm.downstream.custom()
```



```
auto *Intr = cast<IntrinsicInst>(I);
assert(Intr->getIntrinsicID() == Intrinsic::not_intrinsic);
```

Example 2: Constant values vs. materialization

```
%1 = add i32 %0, 42
```

vs.

```
%1 = arith.constant 42 : i32
%2 = arith.addi %0, %1 : i32
```

- Value → User → Constant
 - ConstantData
 - ConstantInt, ConstantFP, ...
 - ConstantDataSequential
 - ConstantAggregate
 - ConstantStruct, ConstantArray, ConstantVector
 - ConstantExpr
 - GlobalValue
 - GlobalIndirectSymbol
 - GlobalObject
 - Function
 - GlobalVariable
 - BlockAddress
 - DSOLocalEquivalent
 - UndefValue
 - PoisonValue

vs.

(N/A)

Deciding on a Unified Design

```
%1 = add i32 %0, 42
```

vs.

```
%1 = arith.constant 42 : i32  
%2 = arith.addi %0, %1 : i32
```

- Constant equality via object identity
- Easier extensibility
- Debug info on constants
- Multi-threading

What is the right design in the long term?
Where is the (near-term) value?

Eliminating `llvm::Constant`

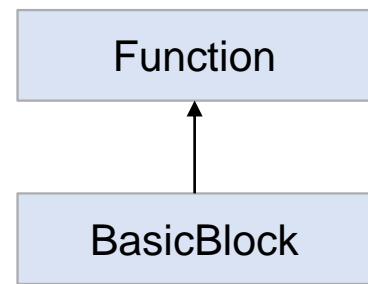
- Value → User → Constant
 - ConstantData
 - ConstantInt, ConstantFP, ...
 - ConstantDataSequential
 - ConstantAggregate
 - ConstantStruct, ConstantArray, ConstantVector
 - ConstantExpr
 - GlobalValue
 - GlobalIndirectSymbol
 - GlobalObject
 - Function
 - GlobalVariable
 - BlockAddress
 - DSOLocalEquivalent
 - UndefValue
 - PoisonValue
- 1. Remove from `GlobalVariable` initializers
 - Standalone representation of constant data
- 2. Remove from metadata
- 3. Introduce symbol reference instruction(s),
detach `GlobalValue` from the Value hierarchy
- 4. Introduce constant instruction(s) for plain data,
remove Constant

Example 3: Regions

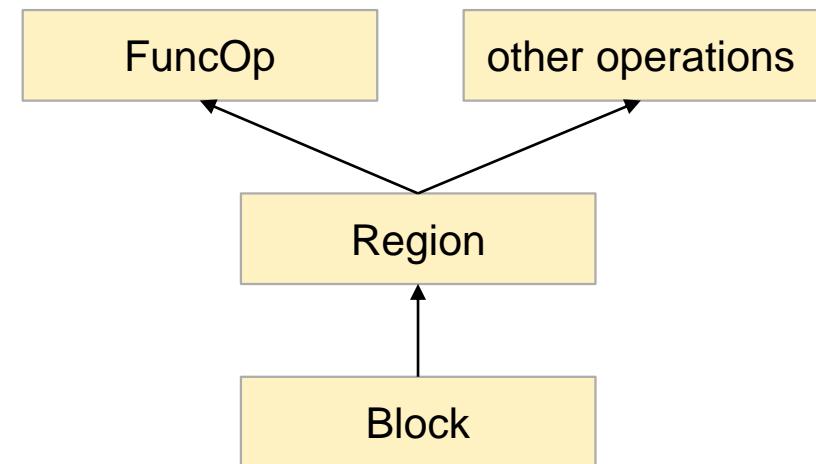
(N/A)

vs.

```
scf.reduce(%cf) : f32 {  
  ^bb0(%lhs: f32, %rhs: f32):  
    %1 = arith.addf %lhs, %rhs : f32  
  scf.reduce.return %1 : f32  
}
```



vs.

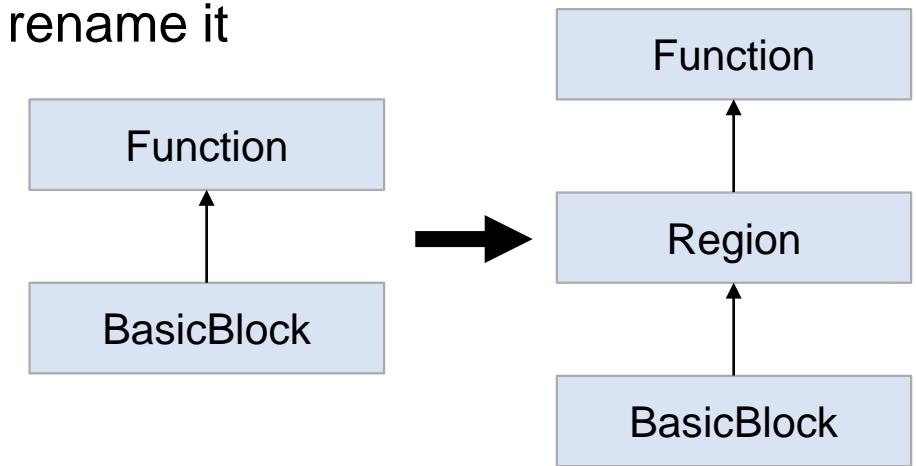


The near-term value of regions in LLVM

- Wave-wide mode in AMDGPU: side-stepping regular control flow
 - Analogous to “unmasked” in ISPC
- Waterfall loops in AMDGPU: highly regular loop with unusual and complex loop condition that does not benefit from generic loop transforms
- Frontends for C-like languages: scopes and lambdas

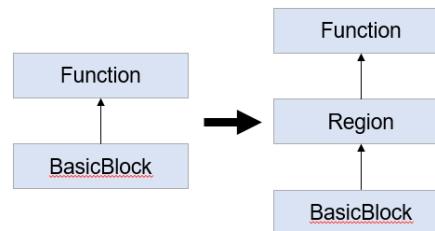
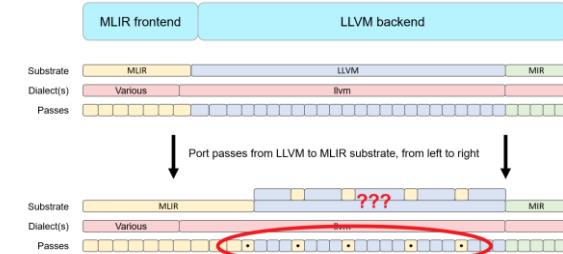
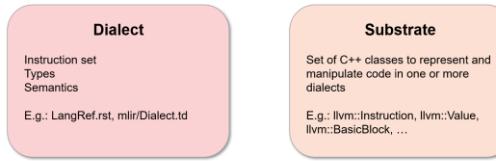
Regions for the LLVM substrate

1. `llvm::Region` exists but means something quite different; rename it
2. Insert Region into the Function/BasicBlock relationship
 - Change `BB->getParent()` to return Region instead of Function; replace by `BB->getFunction()` where necessary and appropriate
 - Iterate over `F->getBody()` instead of `*F`
3. Define a concept of region passes
4. Audit existing function passes: which ones should really be region passes?
5. Add regions to instructions



Recap

- Dialects and substrates
- Benefits of a unified substrate
- Thought experiments on unification of LLVM and MLIR
- The delta matrix
- Feasibility and near-term value of some column refactorings



	Constants	Multiple defined values	Phi nodes vs. B args	Builder interface	Regions	Attributes	Debug info	Metadata	...
Clang									
Flang									
SimplifyRHS									
InstCombine									
Attributor									
Coroutines									
AddressSanitizer									
Atlas analysis									
Target codegen									
Assembly printer & parser									
Bitcode writer & reader									
IR linker									
lldb									
llvm-reduce									
lli									
...									

A vibrant, abstract painting of a tunnel with a bright light at the end. The tunnel walls are composed of dark, textured brushstrokes in shades of black, brown, and green. A bright, white light source at the end of the tunnel creates a strong glow, illuminating the path ahead and casting soft shadows. The overall composition is dynamic and energetic, with a sense of movement and depth.

The End

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Dialect implementation that is generic over the substrate

- Using a generic builder to create a generic instruction.
- Behind the scenes, the template parameter is resolved to the substrate-specific instruction class

```
builder.template create<generic::AddInst>(lhs, rhs);
```

- Using `dyn_cast<T>` with a generic instruction.
- The result type is a function of both the generic target instruction and the substrate-specific operand type, so the result type can be substrate-specific as well.

```
for (auto &inst : basicBlock) {  
    if (auto *alloca = dyn_cast<generic::AllocaInst>(&inst)) {  
        ...  
    }  
}
```